

ESERCITAZIONE 6

NOTAZIONE

Quantizzazione uniforme su $[-V, V]$ con $L=2^b$ livelli.

$$\Delta = \frac{2V}{L}$$

$$P_{eq} \approx \frac{\Delta^2}{12}$$

prob. errore di quantizzazione

Supponendo che le ampiezze del segnale su $[-V, V]$ siano uniformi. $P_x \approx \frac{V^2}{3}$

In decibel

$$P_{dB} = 10 \log_{10}(P)$$

Il rapporto segnale-rumore

$$SNR_{dB} = 10 \log_{10} \left(\frac{P_x}{P_{eq}} \right)$$

ENTROPIA

Entropia di una sorgente discreta:

$$H = - \sum_i p_i \log_2 p_i$$

prob. dei simboli

Esercizio 1

Si consideri un segnale analogico vocale con banda in $[300, 3400]$ (Hz). Il segnale viene campionato e poi quantizzato uniformemente.

Assumiamo $f_s = 8 \text{ kHz}$ (f. campionamento)

Amplitude normalizzata in $[-1, 1]$.

a) $f_s \geq 2B = 2 \cdot 3400 = 6.8 \text{ kHz}$ rispetta la condizione Shannon-Nyquist.

b) $\Delta = \frac{\text{dinamica del segnale}}{L} = \frac{2}{L}$ | $[-1, 1]$ |

Sapendo che $L = 2^b$ # bit di quantizzazione

$$\Delta = \frac{2}{2^b} = 2^{1-b}$$

c) Usando il fatto che $P_{eq} = \frac{\Delta^2}{12}$ determinare il # bit minimo b tale che

$$P_{dB} \leq -20 \text{ dB}$$

$$P_{dB} = 10 \log_{10}(P_{eq}) = 10 \log_{10}\left(\frac{\Delta^2}{12}\right) \leq -20$$

$$\log_{10}\left(\frac{\Delta^2}{12}\right) \leq -2 \Rightarrow \frac{\Delta^2}{12} \leq 10^{-2} \Rightarrow (2^{1-b})^2 \leq 12 \cdot 10^{-2}$$

$$2^{2-2b} \leq \frac{12}{100} \rightarrow 2-2b \leq \log_2 \left(\frac{12}{100} \right)$$

44
-3.06

$$2b \geq 5.06$$

↓
 $b \geq 2.53 \Rightarrow$ # minimo di bit è $\boxed{3}$.

d) Assumendo che il segnale ha cariche uniformi in $[-V, V]$.

$$V=1$$

$$P_x = \frac{V^2}{3} = \boxed{\frac{1}{3}}$$

e) Per $b=3$ e $b=8$ calcolare P_{eq} , SNR_{dB} e bit-rate.

$$\underline{b=3}$$

$$\boxed{P_{eq}} = \frac{\Delta^2}{12} = \frac{2^{2-2b}}{12} = \frac{2^{2-6}}{12} = \frac{1}{2^4} \cdot \frac{1}{12} = \frac{1}{3 \cdot 2^6} = \boxed{\frac{1}{192}}$$

$$\boxed{SNR_{dB}} = 10 \log_{10} \left(\frac{P_x}{P_{eq}} \right) = 10 \log_{10} \left(\frac{\frac{1}{3}}{\frac{1}{192}} \right)$$

bit-rate

$$= 10 \log_{10}(64) = \boxed{18.07 \text{ dB}}$$

$$\boxed{R_b} = f_s \cdot b = 8 \text{ kHz} \cdot 3 \text{ bit} = \boxed{24 \text{ kbit/s}}$$

$$b=8$$

$$\boxed{P_{eq}} = \frac{2^{2-2 \cdot 8}}{12} = \frac{2^{-14}}{12} = \boxed{\frac{1}{3 \cdot 2^{16}}}$$

$$\begin{aligned}
 \boxed{\text{SNR}_{\text{dB}}} &= 10 \log_{10} \left(\frac{P_x}{P_{\text{op}}} \right) = 10 \log_{10} \left(\frac{1}{3} \cdot 8 \cdot 2^{16} \right) \\
 &= 10 \log_{10} (2^{16}) \\
 &= 160 \log_{10} (2) \\
 &\approx 48.17 \text{ dB} \quad \sim 0.30 \text{ M}
 \end{aligned}$$

La differenza tra

$$\begin{array}{c}
 \text{(3hit)} \\
 \text{SNR}
 \end{array}
 -
 \begin{array}{c}
 \text{(3bit)} \\
 \text{SNR}
 \end{array}
 = 48.17 - 18.07 = 30.1 \text{ dB}$$

$$30.1 \text{ dB} = 5 \cdot \boxed{6.02 \text{ dB}} \rightarrow \text{ogni bit migliore di } \underline{6.02 \text{ dB}}$$

f) Dato una finestra $x(n) = \left(\frac{15}{16}, \frac{9}{16}, \frac{3}{16}, -\frac{3}{16}, -\frac{11}{16}, -\frac{15}{16}, \frac{7}{16}, \frac{1}{16} \right)$

Come $b=3$ $L=2^b=8$ livelli

$$x(n) \in [-1, 1] \quad \Delta = \frac{2}{L} = \frac{2}{8} = \boxed{\frac{1}{4}}$$

LIVELLI

$$\left[-1, -\frac{3}{4} \right), \left[-\frac{3}{4}, -\frac{1}{2} \right), \left[-\frac{1}{2}, -\frac{1}{4} \right), \left[-\frac{1}{4}, 0 \right), \left[0, \frac{1}{4} \right), \left[\frac{1}{4}, \frac{1}{2} \right), \left[\frac{1}{2}, \frac{3}{4} \right), \left[\frac{3}{4}, 1 \right)$$



punti medi

$$x_q[n] = \left(\frac{7}{8}, \frac{5}{8}, \frac{1}{8}, -\frac{1}{8}, -\frac{5}{8}, -\frac{7}{8}, -\frac{3}{8}, \frac{1}{8} \right)$$

$$\begin{aligned} P_{x_q} &= \frac{1}{8} \sum_{n=1}^8 x_q^2[n] = \frac{1}{8} \cdot \frac{2 \cdot 7^2 + 1 \cdot 5^2 + 3 \cdot 1^2 + 1 \cdot 3^2}{8^2} \\ &= \frac{1}{8^3} \cdot \left[\underbrace{2 \cdot 49 + 50 + 3 + 9}_{160} \right] \\ &= \frac{160}{8^3} = \frac{20}{64} = \boxed{\frac{5}{16}} \end{aligned}$$

$$e[n] = x_p[n] - x_q[n] = \left(-\frac{1}{10}, \frac{1}{16}, -\frac{1}{16}, \frac{1}{16}, \frac{1}{16}, \frac{1}{16}, \frac{1}{16}, \frac{1}{10} \right)$$

$$P_e = \frac{1}{8} \sum_{n=1}^8 e^2[n] = \frac{1}{8} \cdot 8 \cdot \left(\frac{1}{16} \right)^2 = \left(\frac{1}{16} \right)^2 = \boxed{\frac{1}{256}}$$

$$\begin{aligned} P_x &= \frac{1}{8} \cdot \frac{2 \cdot 15^2 + 4^2 + 2 \cdot 3^2 + 11^2 + 7^2 + 1^2}{16^2} = \frac{1}{8} \cdot \frac{720}{256} = \frac{1}{8} \cdot \frac{45}{16} \\ &= \boxed{\frac{45}{128}} \end{aligned}$$

$$\boxed{SNR_{dB}} = 10 \log_{10} \left(\frac{45/128}{1/256} \right)$$

$$\begin{aligned} &= 10 \log_{10} \left(\frac{45 \cdot 256}{128} \right) = 10 \log_{10}(90) \\ &\approx \boxed{19.54 \text{ dB}} \end{aligned}$$

Quella teorica deve $\approx \underline{18.07 \text{ dB}}$

ESERCIZIO 2

Dato il segnale

$$x(t) = \frac{3}{4} \operatorname{rinc}(t+1) - \frac{1}{8} \operatorname{rinc}(t) - \frac{3}{4} \operatorname{rinc}(t-1)$$

Ricordiamo che

$$\phi_k(t) = \operatorname{rinc}(t-k) \quad k \in \mathbb{Z}$$

base ortogonale
dello spazio segnali a banda $[-\frac{1}{2}, \frac{1}{2}]$

a) Scrivere i coefficienti nell'espansione

$$x(t) = \sum_{k \in \mathbb{Z}} a_k \underbrace{\phi_k(t)}_{\operatorname{rinc}(t-k)}$$

$$a_{-1} = \frac{3}{4}$$

$$a_0 = -\frac{1}{8}$$

$$a_1 = -\frac{3}{4}$$

$$a_k = 0 \quad k \neq \pm 1, 0$$

b) Banda del segnale campionato ($f_s = 1$)

$$\left[-\frac{3}{4}, \frac{3}{4}\right]$$

c) Quarta parte uniformemente $L=3$ livelli.

$$\boxed{\Delta} = \frac{2 \cdot \frac{3}{9}}{3} = \frac{\frac{3}{2}}{3} = \boxed{\frac{1}{2}}$$

$$\left[-\frac{3}{4}, -\frac{1}{4}\right], \left[-\frac{1}{4}, \frac{1}{4}\right], \left[\frac{1}{4}, \frac{3}{4}\right]$$

$\begin{array}{ccc} \text{I}_1 \downarrow & \text{I}_2 \downarrow & \text{I}_3 \downarrow \\ -\frac{1}{2} & 0 & \frac{1}{2} \end{array}$

} punti medi

d) Calcolare coeff. puntiformi \hat{a}_n

$$\begin{array}{l} a_{-1} = \frac{3}{4} \longrightarrow \hat{a}_{-1} = \frac{1}{2} \\ a_0 = -\frac{1}{8} \longrightarrow \hat{a}_0 = 0 \quad a_n = 0 \quad n \neq \pm 1, 0 \\ \hat{a}_1 = -\frac{3}{4} \longrightarrow \hat{a}_1 = -\frac{1}{2} \end{array}$$

$$\begin{aligned} e) \quad \hat{x}_q(t) &= \sum_{n \in \mathcal{Z}} \hat{a}_n \operatorname{rinc}(t-n) \\ &= \frac{1}{2} \operatorname{rinc}(t+1) - \frac{1}{2} \operatorname{rinc}(t-1) \end{aligned}$$

f) Calcolare l'errore di quantizzazione

$$\|x - \hat{x}\|^2 = \sum_{n \in \mathbb{Z}} |e_n - \hat{e}_n|^2 \rightarrow \text{discreto}$$

continuo Persevel

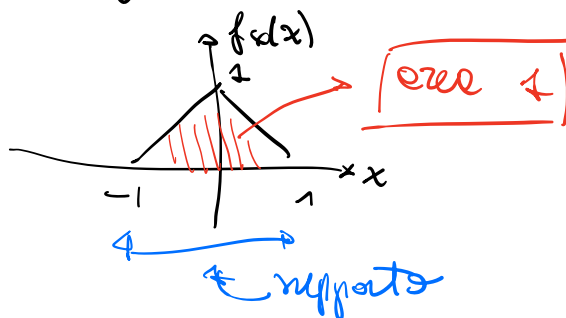
$$\begin{aligned}
 &= (e_{-1} - \hat{e}_{-1})^2 + (e_0 - \hat{e}_0)^2 + (e_1 - \hat{e}_1)^2 \\
 &= \left(\frac{3}{4} - \frac{1}{2}\right)^2 + \left(-\frac{1}{8} - 0\right)^2 + \left(-\frac{3}{4} - \left(-\frac{1}{2}\right)\right)^2 \\
 &= 2 \cdot \left(\frac{1}{4}\right)^2 + \left(\frac{1}{8}\right)^2 = 2 \cdot \frac{1}{16} + \frac{1}{64} = \boxed{\frac{9}{64}}
 \end{aligned}$$

$$\|x - \hat{x}\| = \boxed{\frac{3}{8}}$$

Esercizio 3

Si è X una variabile casuale con densità triangolare

$$f_X(x) = \text{tri}(x)$$



Si consideri X uniformemente come quant. $L=8$
 in $[-1, 1]$.

a) Trovare il passo di quantizzazione

$$\Delta = \frac{2}{8} = \frac{1}{4}$$

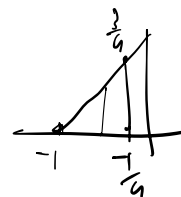
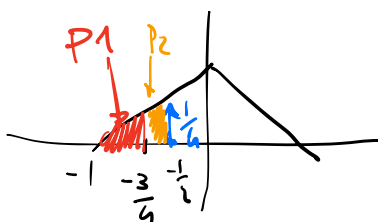
$$\left[-1, -\frac{3}{4}\right), \left[-\frac{3}{4}, -\frac{1}{2}\right), \left[-\frac{1}{2}, -\frac{1}{4}\right), \left[-\frac{1}{4}, 0\right), \left[0, \frac{1}{4}\right), \left[\frac{1}{4}, \frac{1}{2}\right), \left[\frac{1}{2}, \frac{3}{4}\right), \left[\frac{3}{4}, 1\right)$$

$\underbrace{\hspace{1.5cm}}_{p_1} \quad \underbrace{\hspace{1.5cm}}_{p_2} \quad \underbrace{\hspace{1.5cm}}_{p_3} \quad \underbrace{\hspace{1.5cm}}_{p_4} \quad \underbrace{\hspace{1.5cm}}_{p_5} \quad \underbrace{\hspace{1.5cm}}_{p_6} \quad \underbrace{\hspace{1.5cm}}_{p_7} \quad \underbrace{\hspace{1.5cm}}_{p_8}$

b) Calcolare le prob. di ciascun simbolo in uscita del quantizzatore

$$p_1 = \text{vedere } p_1 = \int_{-1}^{-\frac{3}{4}} (1-x) dx = \left[x - \frac{x^2}{2}\right]_{-1}^{-\frac{3}{4}}$$

$$= \frac{1}{4} \cdot \frac{1}{4} \cdot \frac{1}{2} = \frac{1}{32}$$



$$p_2 = \frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2} - \frac{1}{32} = \frac{1}{8} - \frac{1}{32} = \frac{3}{32}$$

$$p_3 = \frac{1}{2} \cdot \frac{3}{4} \cdot \frac{3}{4} - \frac{1}{8} = \frac{9}{32} - \frac{1}{8} = \frac{5}{32}$$

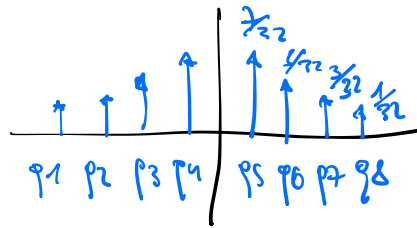
$$p_4 = \frac{1}{2} - \frac{9}{32} = \frac{16}{32} - \frac{9}{32} = \frac{7}{32}$$

$$p_5 = p_6 = \frac{7}{32}$$

$$p_6 = p_3 = \frac{5}{32}$$

$$p_7 = p_2 = \frac{3}{32}$$

$$p_8 = p_1 = \frac{1}{32}$$



c) Calcolo potenza dell'errore di quantizzazione

$$P_{eq} = \frac{\Delta^2}{12} = \left(\frac{1}{4}\right)^2 \cdot \frac{1}{12} = \frac{1}{16} \cdot \frac{1}{12} = \boxed{\frac{1}{192}}$$

$$\Delta = \frac{1}{4}$$

d) Calcolo l'entropia della sorgente quantizzata

$$H = - \sum_i p_i \log_2 p_i = - \frac{1}{32} \cdot \log_2 \frac{1}{32} + \dots$$

$$= \frac{1}{32} \cdot 2 \cdot \log_2 32 + \frac{3}{32} \cdot 2 \cdot \log_2 \left(\frac{32}{5}\right) +$$

$$\frac{5}{32} \cdot 2 \cdot \log_2 \left(\frac{32}{3}\right) + \frac{7}{32} \cdot 2 \cdot \log_2 \left(\frac{32}{7}\right)$$

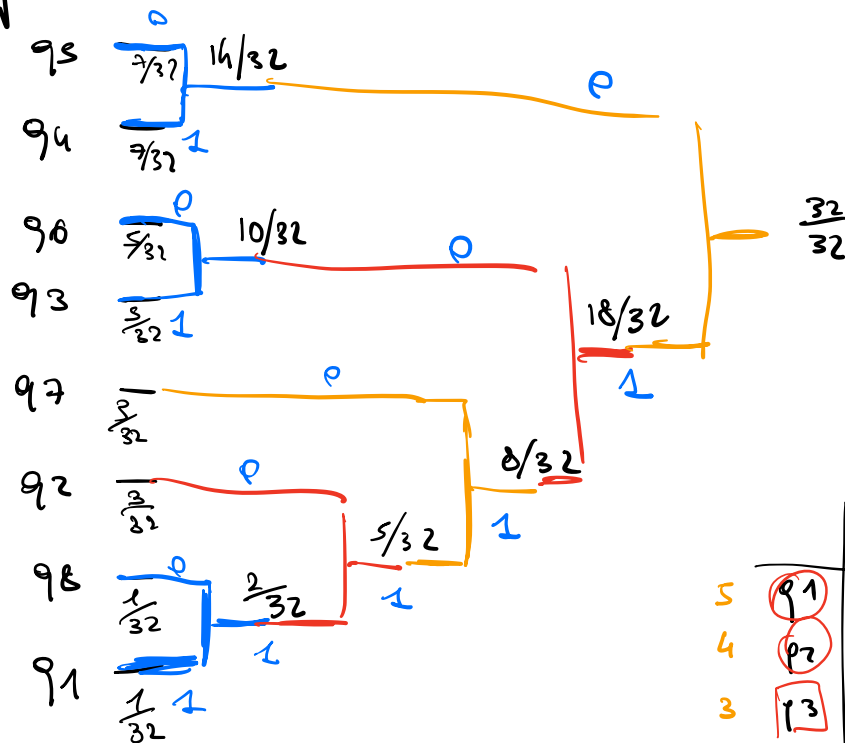
$$\approx 2.749 \text{ bit/simbolo}$$

e) Costruire un possibile codice di Huffman e calcolare la lunghezza media.

$$L = \sum_i l_i \cdot p_i$$

lunghezza codice simbolo i-esimo

Huffman



		Huffman
5	q1	1 1 1 1 1
4	q2	1 1 1 0
3	q3	1 0 1
2	q4	0 1
2	q5	0 0
3	q6	1 0 0
3	q7	1 1 0
5	q8	1 1 1 1 0

$$L = 2 \cdot 5 \cdot \frac{1}{32} + 4 \cdot 4 \cdot \frac{3}{32} + 3 \cdot 1 \cdot \frac{3}{32} + 2 \cdot 3 \cdot \frac{5}{32} + 2 \cdot 2 \cdot \frac{7}{32}$$

q_3, q_6
 q_4, q_5

$$\bar{L} = \frac{10}{32} + \frac{12}{32} + \frac{9}{32} + \frac{30}{32} + \frac{28}{32} = \frac{1}{32} [89]$$

$$\approx 2.781 \text{ bit/simbolo}$$

$$H \leq \bar{L} < H+1$$

verifica
la relazione

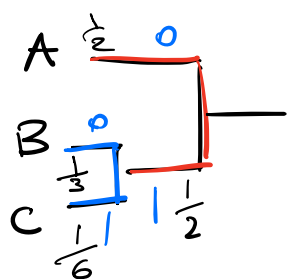
Esercizio 4

Dato una sorgente discreta che emette simboli (senza memoria).

$$X = \{A, B, C\} \rightarrow \text{alfabeto}$$

$$P(A) = \frac{1}{2}, P(B) = \frac{1}{3}, P(C) = \frac{1}{6}$$

a) Costruire Huffman e calcolo \bar{L}_1 lunghezza media



	Huff
A	0
B	10
C	11

cod. prefisso

$$\bar{L}_1 = \overbrace{1 \cdot \frac{1}{2}}^A + \overbrace{2 \cdot \frac{1}{3}}^B + \overbrace{2 \cdot \frac{1}{6}}^C = \frac{1}{2} + \frac{2}{3} + \frac{1}{3} = 1 + \frac{1}{2} = \underline{\underline{\frac{3}{2}}}$$

↓ e.s bit/simboli blocco singolo simbolo

b) Supponiamo ora la sorgente e blocchi di 2 simboli.

$$x_1 x_2 \in \mathcal{X}^2$$

$$\mathcal{X}^2 = \left\{ AA, AB, AC, BA, BB, BC, CA, CB, CC \right\}$$

9 blocchi

c) Calcolare l'entropia della sorgente originale e quella a blocchi di 2.

originale

$$H = -\frac{1}{2} \log_2 \frac{1}{2} - \frac{1}{3} \log_2 \frac{1}{3} - \frac{1}{6} \log_2 \frac{1}{6} \approx 1.459 \text{ bit/simbolo}$$

blocco 2 simboli

iid

$$P(x_1 x_2) = P(x_1) \cdot P(x_2)$$

$$P(AA) = P(A) P(A) = \frac{1}{4}$$

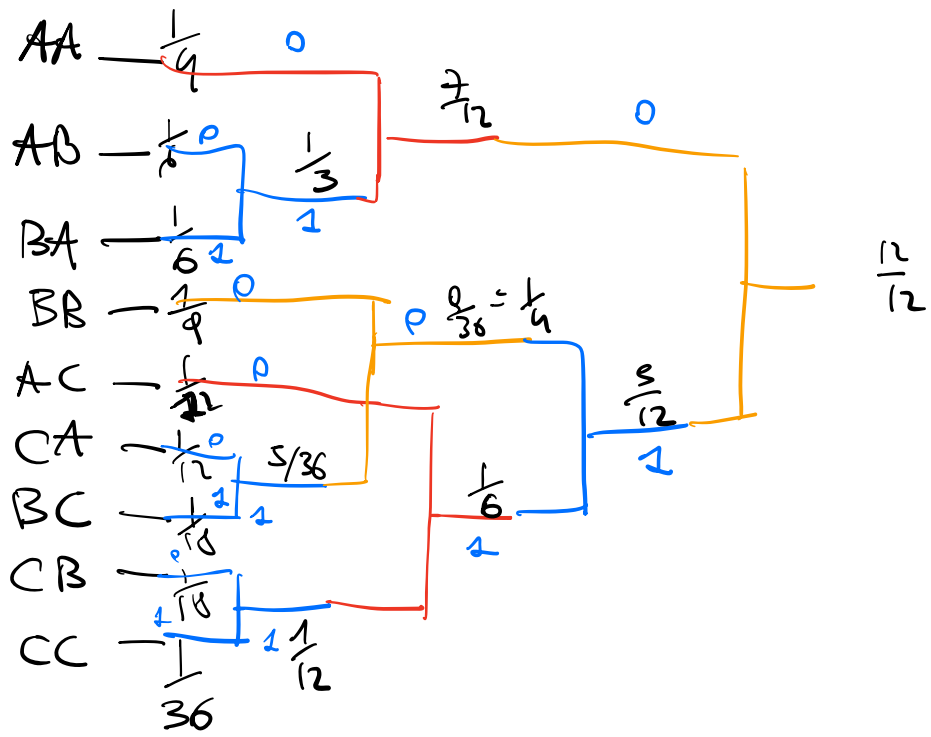
$$H' = 2H \approx 2.918 \text{ bit/blocco}$$

↑ entropie singolo simbolo

entropie

blocco 2 simboli

d) Codice Huffman blocco simbolici



	Huffman
• AA	00
• AB	010
• AC	110
• BA	011
• BB	100
• BC	1011
• CA	1010
• CB	1110
• CC	1111

$$\begin{aligned}
 L_2 &= \overbrace{2 \cdot \frac{1}{4}}^{AA} + \overbrace{2 \cdot 3 \cdot \frac{1}{6}}^{AB, BA} + \\
 &\quad \overbrace{3 \cdot \frac{1}{9}}^{BB} + \overbrace{3 \cdot \frac{1}{12}}^{AC} + \overbrace{4 \cdot \frac{1}{12}}^{CA} + \\
 &\quad \overbrace{2 \cdot 4 \cdot \frac{1}{18}}^{BC, CB} + \overbrace{4 \cdot \frac{1}{36}}^{CC} \\
 &= \frac{1}{2} + 1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{3} + \frac{4}{9} + \frac{1}{9} \\
 &= 1 + 1 + \frac{5}{12} + \frac{5}{9} = \frac{72 + 15 + 20}{36}
 \end{aligned}$$

$$= \frac{107}{36} \approx 2.972 \text{ bit/blocco}$$

$$\downarrow$$
$$\frac{\overline{L_2}}{2}$$

$$= \frac{2.972}{2} \approx 1.486 \text{ bit/simbolo}$$

$$< \overline{L_1} = 1.5 \text{ bit/simbolo}$$

Oss. Codificare blocchi da 2 è più conveniente
in termini di bit/simbolo.

Si può pensare di considerare blocchi di n
simboli $\frac{\overline{L_n}}{n} \xrightarrow{n \rightarrow \infty} H$ (entropia sorgente)